Networking Basics

a moved and

Wolfgang Tremmel academy@de-cix.net

自己的法律法律

Where networks meet

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www.de-cix.net

DE-CIX Management GmbH | Lindleystr. 12 | 60314 Frankfurt | Germany Phone + 49 69 1730 902 0 | sales@de-cix.net | www.de-cix.net





Networking Basics 01 - Networks, Packets, and Protocols

Wolfgang Tremmel academy@de-cix.net

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Networking Basics DE-CIX Academy



- 02 Ethernet
 - 02a VLANs
- 03 The Internet Protocol (IP)

 - 03b Global IP routing
- 04a UDP
- 04b TCP
- 04c ICMP



03a - IP addresses, prefixes, and routing

05 - Unicast, Broadcast, Multicast, Anycast





Networks

- You all know and use networks
- I am not talking about the Internet
- Networks are everywhere
- Example: The road network
- It connects cities, using roads and cars





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- You all know and use networks
- I am not talking about the Internet
- Networks are everywhere
- Example: The railway network
- It connects railway stations, using rails and trains





Attribution: Hbf878, OpenStreetMap contributors https://commons.wikimedia.org/wiki/File:Bahnstrecken_Deutschland_Karte.svg



- You all know and use networks
- I am not talking about the Internet
- Networks are everywhere
- Example: The electrical grid
- It connects producers and consumers of electricity





Attribution: Wolfgang Tremmel

- You all know and use networks
- I am not talking about the Internet
- Networks are everywhere
- Example: The postal network
- It connects senders and receivers of letters and **Packets**















Packets

Just imagine... Using the postal network

- Lets say you have seen this garden shed in a catalog
- And want to mail order it
- It does not fit into one package
- So the sender dismantles it





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Just imagine... Using the postal network

- So the sender dismantles it
- Packs it into several packages
- And sends them to you
- You unpack the packages
- And put the garden shed together again





Attribution: Concord https://commons.wikimedia.org/wiki/File:HGW_Gewaechshaeuser_5_2014_017.JPG



Using a computer network Data is sent in packets

- Computer networks (and the Internet) work in the same way
- Data is most of the time too big to sent as a whole
- So it is cut into packets by the sender
- And put together by the receiver
- Network Protocols take care that the transmission is successful











Protocols

What is a "Protocol"?

- If you want to communicate, you need to speak a common language
- Otherwise you will not understand each other





What is a "Protocol"?

- If you want to communicate, you need to speak a common language
- Otherwise you will not understand each other
- The same is true for computers or other network devices





Protocols used Internet

- The protocol used for the Internet is called "IP"
- IP stands for "Internet Protocol"
 - actually there are two variants of this protocol used
 - IPv4 and IPv6
 - There will be a webinar about IP!





Attribution: cable data by Greg Mahlknecht, map by Openstreetmap contributors https://commons.wikimedia.org/wiki/File:Submarine_cable_map_umap.png





Protocols used Local networks

- The protocol used for your local network at home is called "Ethernet"
- But you also use IP!
- Confused?
 - Multiple protocols are used
 - This is called a Protocol Stack







Protocol Stack OSI reference model



The Protocol Stack **OSI** reference model

- OSI = Open Systems Interconnection
 - Project at ISO (International Organization) for Standardization)
- Defined in the 1970s and 1980s
- Intended as a vendor-independent "real" network model
 - But never took off



Still relevant as reference for teaching

Layer	Name
7	Application
6	Presentation
5	Session
4	Transport
3	Network
2	Data Link
1	Physical



OSI Model Relevance

- It was intended as a working network
- It had some good ideas
 - Like the separation of layers
 - And layers building on each other
- While OSI was writing papers, the Internet guys were already implementing



Layer	Name
7	Application
6	Presentation
5	Session
4	Transport
3	Network
2	Data Link
1	Physical



OSI Model

	Layer	Name
	7	Application
	6	Presentation
	5	Session
	4	Transport
	3	Network
	2	Data Link
DECIX	1	Physical

Internet Model

Layer	Name
5	Application
4	Transport
3	Internet
2	Link
1	Physical

Internet Model Also has layers

- The IP protocols do not exactly fit into the OSI model
- There is also no intention of compliance to OSI
- Internet Protocols care more about architecture and optimization than about layering
- "Running code" most important



Layer	Name
5	Application
4	Transport
3	Internet
2	Link
1	Physical



Internet Model Physical Layer

- Data units are bits or symbols
- via a physical medium
- Like light pulses or electrical signals
- Examples:
 - an optical transceiver
 - an electrical Ethernet port



Layer	Nam
5	Applica
4	Transp
3	Interr
2	Link
1	Physic



Internet Model Link Layer

- Data units are called "Frames"
- Provides node-to-node data transfer
- Examples:
 - Point-to-Point Protocol (PPP)
 - Ethernet



Layer	Nam
5	Applica
4	Transp
3	Interr
2	Linl
1	Physi



Internet Model IP / Internet Layer

- Data units are called "Packets"
- Provides source to destination transport
- Needs addresses for hosts
- Examples:
 - IPv4
 - IPv6



Layer	Nam
5	Applica
4	Transp
3	Interr
2	Linl
1	Physi



Internet Model **Transport Layer**

- May provide flow control, reliability, congestion avoidance
- Not all of them in all protocols
- Also contains information about the next layer up
- Examples:
 - TCP (flow control, reliability, congestion) avoidance)



UDP (none of the above)

Layer	Nam
5	Applica
4	Transp
3	Interr
2	Linl
1	Physi



Internet Model **Application Layer**

- no special name for data units, just "Data"
- Contains all application protocols
- Examples:
 - SMTP, HTTP, SSH and all others



Layer	Nan
5	Applica
4	Trans
3	Inter
2	Lin
1	Phys



Conclusion



Conclusion This is what you should remember

- Networks are everywhere
- Data is sent in **packets** via a network
- other
- The OSI model defines a network as a number of layers
- The Internet does not exactly fit into the OSI model



• A protocol is a common language devices speak so they understand each

Networking Basics

02 - Ethernet

Wolfgang Tremmel academy@de-cix.net

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01 - Networks, Packets, and Protocols

03a - IP addresses, prefixes, and routing

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Ethernet

A Modern Ethernet Device Nokia 7950

- As used by DE-CIX
- Connects 100s of devices
- using optical interfaces
- with speeds up to 400Gbps







Another Modern Ethernet Device Fritzbox

- as used at home
- connects 4 devices directly
- using copper interfaces
- with speeds up to 1Gbps





Attribution: Wolfgang Tremmel

So why does the symbolic drawing of Ethernet look like this?









It began in Hawaii: ALOHA-Net



#
ALOHA-Net University of Hawaii, 1971

- Radio based network
- To interconnect sites
- Simple principle:
 - If you have data to send, send it
 - If you receive something while sending, stop and try again later







Robert Metcalfe - Xerox PARC



Ethernet Xerox PARC, 1973

- Instead of radio, uses a coax cable
 - For higher bandwidth
 - And more reliability
- Inspired by ALOHAnet
- Standardized in 1980
- Ethernet II in 1982, standardized as IEEE 802.3 in 1983 DECIX



https://commons.wikimedia.org/wiki/File:Parcentrance.jpg

10BASE5 10 Mbit/s Ethernet

- 10 Mbit/s
- BASE uses baseband modulation
- 5 500m max. segment length
- Hardware:
 - 1cm thick coax cable
 - "Vampire-Tap" Transceivers





Attribution: Robert.Harker at English Wikipedia https://commons.wikimedia.org/wiki/File:10Base5transcievers.jpg





10Base5 Ethernet Remember the drawing





10Base5 Ethernet Remember the drawing





10Base5 Ethernet Remember the drawing







10BASE2 still only 10 Mbit/s Ethernet

- Hardware:
 - thin coax cable
 - BNC-"T"-connectors
- Up to 200m total length
- "Cheapernet"
- mid to late 1980s





https://twitter.com/the_mutax/status/1303700688745226240



Attribution: Dmitry Nosachev https://commons.wikimedia.org/wiki/File:3Com_3C509BC_Ethernet_NIC.jpg 44



10Base-T still only 10 Mbit/s Ethernet

- Hardware:
 - two pairs of twisted copper wires
 - 8P8C (RJ45) plastic connector
- Since 1988
- Needs an active component (hub) or switch) to interconnect







Attribution: Dmitry Nosachev https://commons.wikimedia.org/wiki/File:3Com_3C509BC_Ethernet_NIC.jpg 45



Competing standards



Token Ring 1984 - 1990s

- Developed by IBM
- 4Mbit/s, later 16Mbit/s
- Deterministic access
- Needs central Multistation Access Unit
- More complex than Ethernet
- More expensive than Ethernet **DE CIX**

a)





FDDI late 1980s - 1990s

- Fiber Distributed Data Interface
- Optical network
- 100Mbit/s speed, up to 200km size
- Frame-size of 4352 bytes
- double ring topology
- made obsolete by GigabitEthernet





Attribution: Maximilian Wilhelm



Attribution: <u>Vincent van der kussen</u> at <u>nl.wikibooks</u> https://commons.wikimedia.org/wiki/File:FDDI_Concentrator.jpeg







Back to Ethernet



Ethernet is a *broadcast* network where all devices are connected to a *shared* medium



Broadcast network One is sending, everybody is receiving

- All stations share one medium
- Only one station at a time can send data
- If two stations start sending at the same time, a collision occurs
 - Both stop sending, wait for a random time, then retry
 - This was one of the main criticisms (no guaranteed delivery)



Broadcast network One is sending, everybody is receiving

- Everybody is receiving everything
- How to avoid overload / unnecessary processing of data?
 - Each station has a unique 48-Bit address
 - Receivers address is at the beginning of each frame
 - And can be processed by the network card
 - Only frames with matching address or broadcast frames are forwarded to the CPU



Ethernet Frame Structure

Preamble	SF D	Destination MAC Address	Source MAC Address	Ethertype	Payload	Checksu
10101010 10101010 10101010 10101010	10101011	48 Bits 6 Octets	48 Bits 6 Octets	16 Bits 2 Octets	46 - 1500 Octets	32 Bits 4 Octets





Ethernet Frame Structure



Prea	mble	SF D	Destination MAC Address	Source MAC Address	Ethertype	Payload	Checksu
10101010 10101010	10101010 10101010	10101011	48 Bits 6 Octets	48 Bits 6 Octets	16 Bits 2 Octets	46 - 1500 Octets	32 Bits 4 Octets

- Preamble 56 bits of 10101010....
- Start of frame marker 8 bits: 10101011
- Destination MAC address
- Source MAC address
- EtherType (or length)
- Payload
- 32 bit checksum



Ethernet Addressing



• 48 Bit address - 6 octets

281 trillion possible addresses

managed by IEEE

- you can purchase blocks of addressesnotation examples:
 - 00:26:b0:d8:3d:8a
 - 0026.b0d8.3d8a
 - 00-26-b0-d8-3c-8a

MAC Address	Ethertype	Payload	Checksu
48 Bits	16 Bits	46 - 1500 Octets	32 Bits
6 Octets	2 Octets		4 Octets



Ethernet Addressing



- Two bits in first octect have special meaning
- one for local vs. globally unique addresses
 - unique: usually "burned" into the hardware by manufacturer
- one for unicast vs. multicast







EthernetSpecial Addresses



Preamble SF D	Destination MAC Address	Source MAC Address	Ethertype	Payload	Checksu
10101010 10101010 10101010 10101010 10101011		48 Bits 6 Octets	16 Bits 2 Octets	46 - 1500 Octets	32 Bits 4 Octets

• FF:FF:FF:FF:FF:FF

The *broadcast* address

• Received by all nodes



Ethernet Ethertype

Preamble				SF D	Destination MAC Address	Source
10101010	10101010	10101010	10101010	10101011	48 Bits 6 Octets	

Was once used to indicate size of payload

- Using values up 1500
- \rightarrow Ethertype values start at 1536
- Some well-known values:

0x0800	IPv4
0x86dd	IPv6
0x0806	ARP
0x8100	VLAN Tagged





Ethernet Today



Ethernet connections In data centers

- Usually optical fibres are used
- Various types exist (single mode, multi mode)
- Speeds are 1 GBit/s, 10 GBit/s, 100 GBit/s or 400 GBit/s
- Connections are between a switch and an end device







Ethernet at home 10Base-T

- Only wire-based connections are in use
- Speeds are 100Mbit/s or 1Gbit/s
- With a switch as a center
- Wireless Ethernet WIFI is most common







Ethernet at home 10Base-T

- 10Base-T (twisted pair) requires a central device
- To replace the yellow coax cable
- Early devices: a hub
 - Function: What is received on one port is broadcasted out on all other ports





Attribution: Zac67 https://commons.wikimedia.org/wiki/File:HP_EtherTwist_Hub8.jpg



Ethernet Switch Ethernet today

- Instead of a hub, a switch is common today
- Advantage:
 - a switch learns which devices are connected to which port
 - and only sends frames on ports they are destined to
 - fallback: unknown destinations are still broadcasted on all ports





Attribution: Wolfgang Tremmel

But... Ethernet still....

- ...usually has a max payload size of 1500 octets
 - "jumbo frames" with 9000 octets exist, but are not commonly used
- ...uses 48-bit addresses
- ... is a broadcast medium.
 - but today switches are used and connections are point-to-point



Network layers - Internet Model **Ethernet: Link Layer**

- Data units are called "Frames"
- Provides node-to-node data transfer
- Examples:
 - Point-to-Point Protocol (PPP)
 - Ethernet



Layer	Nam
5	Applica
4	Transp
3	Interr
2	Link
1	Physic



Conclusion



Please remember.... **Facts about Ethernet**

- Ethernet is a **broadcast** network
- It uses **48-Bit** addresses
 - Which are globally **unique**
- Ethernet frames have usually max. 1500 octets payload
- Today switches interconnect devices



Networking Basics 02a - Ethernet + VLANs

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Wolfgang Tremmel academy@de-cix.net

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- U4b TCP
- 04c ICMP



01 - Networks, Packets, and Protocols

05 - Unicast, Broadcast, Multicast, Anycast





Ethernet

A typical Ethernet In an office building or a home

Ethernet is a *broadcast* network where all

devices are connected to a shared medium

 End devices (computers) connected to the switches





Now you want a second network


Now you want a second network

- For guests
- Or your telephones
- Or for network connected "things"



Second network To keep separate things separate

- For example: Guest network
- Duplicate everything?
- No need it's easier
- VLANs to the rescue!







VLANS Virtual LANs



Ethernet Frame

Preamble			е	SF D	Destination MAC Address	Source MAC		
10101010	10101010	10101010	10101010	10101011	48 Bits 6 Octets	48 B 6 Oct		



Ether	'n	et Frame	• Some	well-	know	n valı/ ו	JES: Pv4		
VLAN ta	ago	jed		0x86d	d	I		Pv6	
				0x080	6	ŀ	ARP		
				0x810	0	VLAN	I Tagg	ed	
Preamble	SF D	Destination MAC Address	Source MAC Address	Ethertype	P	Payload		Checksum	
1010101010101010101010101010101	0101011	48 Bits 6 Octets	48 Bits 6 Octets	16 Bits 2 Octets	46-1	500 Octet	ts	32 Bits 4 Octets	
Preamble	SF D	Destination MAC Address	Source MAC Address	VLAN H (801	leader .1Q)	Ethertype	F	Payload	Chec
10	10101011	48 Bits 6 Octets	48 Bits 6 Octets	0x8100	VLAN	16 Bits 2 Octets	42 -	1500 Octets	32 4 O
			77						

Ethe VLAN	Prne tagg	et Frame	• Some	- well 0x0800	know 0 d	n valu I	Jes: Pv4 Pv6		
				0x080	6	ŀ	ARP		
				0x810	0	VLAN	I Tago	ged	
	SF	Destination MAC							
Preamble	D	Address	Source MAC Address	Ethertype		ayload		Checksum	
10	0101010101011	48 Bits 6 Octets	48 Bits 6 Octets	16 Bits 2 Octets	46-1	500 Octe	S	32 Bits 4 Octets	
Preamble	SF D	Destination MAC Address	Source MAC Address	VLAN F (801.	leader 1Q)	Ethertype		Payload	Cheo
10101010101010101010101010	101010101011	48 Bits 6 Octets	48 Bits 6 Octets	0x8100	VLAN	16 Bits 2 Octets	42 -	- 1500 Octets	32 4 O
			77						



Ethernet VLAN tagged frame

F	Preamble	SF D	Destination MAC Address	Source MAC Address	VLAN H (801	Header .1Q)	Ethertype	Payload	Check
10101010	10101010101010101010101	00101011	48 Bits 6 Octets	48 Bits 6 Octets	0x8100	VLAN	16 Bits 2 Octets	42 - 1500 Octets	32 B 4 Oct





VLAN Header (801.1Q)									
ts	3 Bits	1 Bit	12 Bits						
00	Priority	May be drop ped?	VLAN ID 1-4094						



Multiple networks Use VLANs to separate

• You can have multiple VLANs on one physical infrastructure







Multiple networks Use VLANs to separate

- You can have multiple VLANs on one physical infrastructure
- Connections can have one or multiple VLANs on them
- Connections which carry multiple VLANs are called "trunk"







How to set it up?



How to set it up? Building an Ethernet with VLANs

- You remember Ethernet switches?
- You might have one in your basement
- Ethernet switches connect devices to each other
- Ethernet switches also can connect to other switches











Ethernet Switch



Ethernet Switch









Ethernet Switch







Ethernet Switch

Ethernet Switch



























And how does it work?











Guest



VLANs at DE-CIX





VLANS at DE-CIX How we use them

- VLANs can deliver multiple LANs on one trunked port
- Each tagged with a different VLAN ID
- Like we used to separate "Home" and "Guest" network







VLANs at DE-CIX How we use them

- Customers connect to DE-CIX via Ethernet
- Standard connection is a untagged access port







VLANS at DE-CIX How we use them

- Customers connect to DE-CIX via Ethernet
- Standard connection is a untagged access port
- But we can also deliver via a tagged trunk-like port







VLANS at DE-CIX How we use them

- Customers connect to DE-CIX via Ethernet
- Standard connection is a untagged access port
- But we can also deliver via a tagged trunk-like port
- And on a trunk-like port we can deliver multiple services







VLANs at DE-CIX Connect to the Cloud

- The same way we connect customers to (multiple) Cloud service providers
- At DE-CIX the VLAN ID on each end does not have to be the same!






Conclusion



Please remember.... **Facts about VLANs**

- Ethernet is a **broadcast** network
- VLANs set up virtual LANs on a common physical infrastructure
- VLAN IDs run from 1 4094
 - It is recommended to **not use VLAN 1** (if possible)
- DE-CIX uses VLANs for multiple service delivery on one physical port



Networking Basics 03 - The Internet Protocol (IP)

Wolfgang Tremmel academy@de-cix.net

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Protocol

What is a "Protocol"?

- If you want to communicate, you need to speak a common language
- Otherwise you will not understand each other





What is a "Protocol"?

- If you want to communicate, you need to speak a common language
- Otherwise you will not understand each other
- The same is true for computers or other network devices





Protocol Stack Multiple protools building on each other



Internet Model Physical Layer

- Light pulses and electrical signals
- Lasers and fibres
- Electrical cables



Layer	Nam
5	Applica
4	Transp
3	Interr
2	Linl
1	Physi



Attribution: Cjp24 https://commons.wikimedia.org/wiki/File:12_Optical_fibers_(1).jpg



Internet Model Link Layer

- Data units are called "Frames"
- Provides hop-to-hop (node-to-node) data transfer
- Examples:
 - Ethernet



	Layer	Nam
	5	Applica
	4	Transp
	3	Interr
	2	Lin
-	1	Physi



Attribution: Wolfgang Tremmel



Ethernet Frame Structure

Preamble	SF D	Destination MAC Address	Source MAC Address	Ethertype	Payload	Checksu
10101010 10101010 10101010 1010101	0 10101011	48 Bits 6 Octets	48 Bits 6 Octets	16 Bits 2 Octets	46 - 1500 Octets	32 Bits 4 Octets

EtherTypePayload





Ethernet Frame Structure

Ethertype	Payload	Checksur
16 Bits 2 Octets	46 - 1500 Octets	32 Bits 4 Octets





Encapsulation Packets inside packets

- The payload of Ethernet is IP
- Encapsulation is like Russian dolls
- So we have an IP packet inside an Ethernet frame





Attribution: Fanghong. derivative work: Greyhood https://commons.wikimedia.org/wiki/File:Matryoshka_transparent.png



Internet Model IP / Internet Layer

- Data units are called "Packets"
- Provides source to destination (end-to-end) transport
- Needs addresses for entities
- Examples:
 - IPv4
 - IPv6



Layer	Nam
5	Applica
4	Transp
3	Interr
2	Linl
1	Physi



IP - Version 4 (IPv4) Header + Payload

	Ether
IPv4 Header	
20-60 Bytes	



rnet Payload

IP Payload

1440-1480 Bytes

IPv4 Header Some parts to point out

- Starts with version and length
- Total length of packet
- Important: Time to live (TTL)
- Protocol: Type of payload
- Source / Destination address 32 bits
- Options (optional)





IP - Version 6 (IPv6) Header + Payload

	Ethernet Payload
IPv6 Fixed ^{*)} Header	IP Pa
40 Bytes	1460



IP Payload

1460 Bytes



IPv6 Header Looks less complicated!

- Starts with version and some labels
- Payload length in bytes (0-65535)
- Next Header you can chain more headers
 - replaces protocol field
- Hop Limit replaces TTL







IP Addresses - IPv4



IP Addresses IPv4

- 32 bit in length
- you might have heard of Class-1, -12 -C addresses
- there is no such thing anymore!
 - since 1993!

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- all usable IP addresses are equal
- more about this in another webinar

192.0.2.123

 1 1 0 0
 0 0 0 0
 0 0 0 0
 0 0 0 0
 0 0 0 0
 0 0 1 0
 0 1 1 1
 1 0 1 1

 1 2 3 4
 5 6 7 8
 9 10 11 12
 13 14 15 16
 17 18 19 20
 21 22 23 24
 25 25 27 28
 29 30 31 32

IP Addresses IPv4

- 32 bit in length
 - 4.294.967.296 possible addresses
- written as 4 decimal numbers separated by dots "."
- some addresses are reserved / not usable



192.0.2.123

 1 1 0 0
 0 0 0 0
 0 0 0 0
 0 0 0 0
 0 0 0 0
 0 0 1 0
 0 1 1 1
 1 0 1 1

 1 2 3 4
 5 6 7 8
 9 10 11 12
 13 14 15 16
 17 18 19 20
 21 22 23 24
 25 25 27 28
 29 30 31 32

IPv6!

Development started: 1994 First published: 1995

IP Addresses - IPv6



IP Addresses IPv6

- 128 bit in length possible addresses: 340282366920938463463374607431768211456
- there are lots of IPv6 addresses available
- written as hexadecimal numbers separated by colons ":"
 - double-colon "::" means fill up with zeros here
- some addresses are reserved / not usable



2001:db8:274f:400:226:b0ff:fed8:3d8a

2001:db8::1

2001:db8:0:0:0:0:0:1

2001:0db8:0000:0000:0000:0000:0000:0000:1





Internet Protocol How did it all start?



History of IP It started in the 60s

- To debunk a myth:
 - It was not invented to survive a nuclear war!
- But it was funded by DARPA a military research agency
- To connect research facilities to share (computing) resources



ARPANET/MILNET GEOGRAPHIC MAP, APRIL 1984



https://commons.wikimedia.org/wiki/File:ARPANET_-_MILNT_Diagram_1984.jpg



Timeline Some days in the early history of IP

- 1961 concept of packet switching network
- 1967 Plan for "ARPAnet"
- 1969 first ARPAnet node, first RFC published
- 1972 first public demonstration of ARPAnet
- 1974 TCP/IP protocol described, "Internet" first used

1983 - ARPAnet switches from NCP to TCP/IP **DE CIX**



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Why was IP so successful? The role of documentation

- There were other, competing protocols
 - Do you remember BITnet, DECnet? OSI?
 - Either vendor-proprietary or just theory
- IP Protocols were evolving more quickly
- Everything was open and still is
 - 1969 first Request for comments (RFC) published
 - Memos, best practices, standards published as RFC



Today: Well established and open standard for publishing

Updated by: <u>1349</u>, <u>2474</u>, <u>6864</u> RFC: 791

INTERNET STANDARD

INTERNET PROTOCOL

DARPA INTERNET PROGRAM PROTOCOL SPECIFICATION

September 1981

prepared for

Defense Advanced Research Projects Agency Information Processing Techniques Office 1400 Wilson Boulevard Arlington, Virginia 22209

by

Information Sciences Institute University of Southern California 4676 Admiralty Way Marina del Rey, California 90291

September 1981

Internet Protocol

https://commons.wikimedia.org/wiki/File:ARPANET_-_MILNT_Diagram_1984.jpg



Conclusion



Things you should remember The IP Protocol(s)

- Internet Protocol (IP) takes care of end-to-end communication
- IPv4 and IPv6 coexist
- IP packets consist of header and payload
- IPv4 and IPv6 headers are different
 - But both contain source- and destination addresses
 - IPv4 addresses are 32 bit long, IPv6 addresses are 128 bit long



Payload can yet be another protocol

Networking Basics 03a - IP: Addresses, Prefixes and Routing

Wolfgang Tremmel academy@de-cix.net

目前加強相關的目的

Where networks meet

DECIX

GERMAN NETWORK OPERATORS GROUP



www.de-cix.net

DE-CIX Management GmbH | Lindleystr. 12 | 60314 Frankfurt | Germany Phone + 49 69 1730 902 0 | sales@de-cix.net | www.de-cix.net





Networking Basics DE-CIX Academy

- 02 Ethernet
 - 02a Ethernet + VLANs
- 03 The Internet Protocol (IP)
 - 03b Global IP routing
 - 03a IP addresses, prefixes, and routing
- 04a UDP
- 04b TCP
- 04c ICMP





01 - Networks, Packets, and Protocols

05 - Unicast, Broadcast, Multicast, Anycast

IP - the Internet Protocol



Internet Model IP / Internet Layer

- Data units are called "Packets"
- Provides source to destination (end-to-end) transport
- Needs addresses for entities



Layer	Nam
5	Applica
4	Transp
3	Interr
2	Lin
1	Physi



IPv4 Addresses 32 bit long

- 32 bit in length
 - 4.294.967.296 possible addresses
- written as 4 decimal numbers separated by dots "."
- some addresses are reserved / not usable



192.0.2.123

1100 0000 0000 0000 0000 0010 0111 1011

IPv6

Development started: 1994 First published: 1995



IPv6 Addresses 128 bit long

- 128 bit in length possible addresses: 340282366920938463463374607431768211456
- there are lots of IPv6 addresses available
- written as hexadecimal numbers separated by colons ":"
 - double-colon "::" means fill up with zeros here
- some addresses are reserved / not usable

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2001:db8::1

2001:0db8:0000:0000:0000:0000:0000:0000:1

0 01 02 03 04 05 06 07 08 09 0a 0b 0c 0d 0e 0f 10 11 12 13 14 15 16 17 18 19 1a 1b 1c 1d 1e 1f 20 21 22 32 425 26 27 28 29 2a 2b 2c 2d 2e 2f 30 31 32 33 44 55 46 47 48 49 4a 4b 4c 4d 4e 4f 50 51 52 53 54 55 56 57 58 59 5a 5b 5c 5d 5e 5f 60 61 62 63 64 65 66 67 68 69 6a 6b 6c 6d 6e 6f 70 71 72 73 74 75 76 77 78 79 7a 7b 7c 7d 7e 7f 7a 7b 7c 7d 7e 7f



Internet Model IP / Internet Layer

- Data units are called "Packets"
- Needs addresses for entities
- Provides source to destination transport
 - End-to-End Transport



Layer	Nam
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End-to-End Transport















Intermission: Language



Routing

Routing

"A method of finding paths from origins to destinations in a network such as the Internet, along which information can be passed."





https://en.wiktionary.org/wiki/routing



Routing

"A method of finding paths from origins to destinations in a network such as the Internet, along which information can be passed."

root



https://en.wiktionary.org/wiki/routing



rooting "A hole formed by a pig when it roots in the ground"

https://en.wiktionary.org/wiki/rooting







Routing How a *router* works

- All IP packets have a destination IP address
- Depending on the destination IP address a next hop is chosen
- For this, each router has a large lookup-table
 - This is called the *routing table*
- It contains not single IP addresses, but *Prefixes*





IP Prefixes



IPv6 - Addresses



2003:de:274f:400:226:b0ff:fed8:3d8a

IPv6 - Refresses

2003:de:274f:400:206:b0ff:fed8:3d8a



3c 3d 3e 3f 40 41 42 43 44 45 46 47 48 49 4a 4b 4c 4d 4e 4f 50 51 52 53 54 55 56 57 58 59 5a 5b 5c 5d 5e 5f 60 61 62 63 64 65 66 67 68 69 6a 6b 6c 6d 6e 6f 70 71 72 73 74 75 76 77

128 Bits long



IPv6 - Prefixes

2003:de:274f:400::/64



Prefix-Length: 0-128

Host-part all zero

128 Bits long



IPv4 - Addresses

10.3.8.17







10.3.8.0722

15 16 17 18 19 20 21 22 23 24 25 25 27 28 29 30 31 32

32 Bits long







IP Addresses and Prefixes Prefix or not?

	IPv4	IPv6
Length	32 Bit	128 Bit
	0-32 Prefix Length	0-128 Prefix Length
Notation	4 Numbers, 0-255	8 Numbers, 0-fffff
Separator		
Prefix: Host part	all zero	
Address: Host part	not all zero / not all one	
Example (Prefix)	198.51.100.0/24	2001:db8:4f30::/48



















<u>To: 198.51.100.17</u>







<u>To: 198.51.100.17</u>







<u>To: 198.51.100.17</u>

Apply the netmask: /24

/24 = 24 bits in network part 11111111111111111111100000000

255.255.255.0 "bitwise logical and" with IP address

¥

198.051.100.017 255.255.255.000

198.051.100.000







<u>To: 198.51.100.17</u>

Apply the netmask: /24

/24 = 24 bits in network part 1111111111111111111111000000000

255.255.255.0 "bitwise logical and" with IP address

198.051.1**00.017** 255.255.25**5.000**

=

198.051.100.000 =







- You now know how routing works
 - The router has a routing table with IP prefixes
 - The destination address is used to select a best matching prefix
 - The routing table tells the router the "next hop"

But how does the information get into the routing table?



IP Routing But how does the information get into the routing table?

- Someone types it in
 - This is called "static routing"
 - Simple, often used, but does not scale
- Routers "talk" to each other
 - This is called "dynamic routing"
 - And the protocols used are called "routing protocols"



Examples of such protocols are **BGP**, OSPF, IS-IS, RIP, EIGRP

Conclusion



Conclusion We introduced a lot of new terms in this webinar

- "Routing"
 - which information can be passed." https://en.wiktionary.org/wiki/routing
- "Router"
 - A device which routes
- "IP Prefix"
 - A network address and a prefix length
- "Routing Table"



• "A method of finding paths from origins to destinations in a network such as the Internet, along

Networking Basics 03b - Global IP routing

Wolfgang Tremmel academy@de-cix.net

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5	Applica
4	Transp
3	Interr
2	Linl
1	Physi



IPv6 - Prefixes

2003:de:274f:400::/64



Prefix-Length: 0-128

Host-part all zero

128 Bits long




















How did it start?



Internet in the 1990s Also some personal history

- 2Mbit/s was a "fast connection"
- Backbone networks were 34Mbit/s
- End customers connected via ISDN (modern) or telephone modems
- Standard end customer speed was 64kBit/s



Success rate is 99 percent (999999/100000), round-trip min/avg/max

Hssi12/0 is up, line protocol is up

Hardware 15 CXBus Haal

Internet address is 174.122.225.1/30

MTU 1500 bytes, BW 34100 Kbit, DLY 200 usec, rely 255/255, 10 Encapsulation HDLC, Loopback not set, keepalive set (10 sec) Last input 00:00:02, output 00:00:00, output hang never Last clearing of "show interface" counters never Output queue 0/40, 0 drops; input queue 0/75, 0 drops 5 minute input rate 7051000 bits/sec, 622 packets/sec 5 minute output rate 7052000 bits/sec, 623 packets/sec 1449469 packets input, 2033725506 bytes, 0 no buffer Received 304 broadcasts, 0 runts, 0 giants

) parity

1 input errors, 1 CRC, 0 frame, 0 overrun, 0 ignored, 0 about 1449504 packets output, 2033727908 bytes, 0 underruns 0 output errors, 0 applique, 1 interface resets 0 output buffer failures, 0 output buffers swapped out 1 carrier transitions

gw0



Internet in the 1990s Situation in Germany

- In the early 1990s there were three commercial ISPs in Germany
- Each had a connection to the US
- So traffic had to go via the Atlantic twice in worst case
- And was expensive





Internet in the 1990s The idea

- Idea: Exchange traffic directly
- Via cables?
 - Each of this early ISPs was in a different city
 - Cables / leased lines were expensive
 - Meet in the middle!









Internet in the 1990s The solution

- Meet in the middle!
 - DE-CIX was born
- Direct connection between commercial providers
- To exchange traffic
 - Only their own and their own customers



This is Peering!







Peering Noun

peering (plural peerings)

1. The act of one who peers; a looking around.

usually without charge or payment.

https://en.wiktionary.org/wiki/peering



2. (Internet) The act of carrying communications traffic terminating on one's own network on an equivalency basis to and from another network,

Peering A typical Internet Service Provider

















Peering More direct via peering

17 27

X







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シアド

X





Peering Hierarchy Peering on multiple levels

- Peering happens usually between equal size networks
- Peering takes place on all network levels
- The "top ones" only peer with each other
 - They are called "Tier-1 networks"









Why Peering?



Why Peering? Historical view

- In the early days, peering was cheaper than routing traffic via upstream providers
- Also, it kept traffic from crossing the Atlantic twice



s rate is 99 percent (99999/100000), round-trip min/avg/max = 16.

0 is up, line protocol is up are is cxBus HSSI net address is 194.122.225.1/30 500 bytes, BW 34100 Kbit, DLY 200 usec, rely 255/255, load 52/2 sulation HDLC, loopback not set, keepalive set (10 sec) input 00:00:02, output 00:00:00, output hang never clearing of "show interface" counters never relearing of "show interface" counters never aqueue 0/40, 0 drops; input queue 0/75, 0 drops te input rate 7051000 bits/sec, 622 packets/sec te output rate 7052000 bits/sec, 623 packets/sec te output rate 7052000 bits/sec, 623 packets/sec 9469 packets input, 2033725506 bytes, 0 no buffer eived 304 broadcasts, 0 runts, 0 giants 0 parity nput errors, 1 CRC, 0 frame, 0 overrun, 0 ignored, 0 abort 9504 packets output, 2033727908 bytes, 0 underruns

utput errors, O applique, 1 interface resets utput buffer failures, O output buffers swapped out arrier transitions



Why Peering? Peering today

- Upstream prices are down
- But peering is still important
 - Keeping local traffic local
 - Reducing latency
 - Keep your traffic under control
- At an Internet Exchange (IXP) you can peer with multiple networks via one common infrastructure





Conclusion



Conclusion Global IP routing

- Networks (Internet Service Providers) announce IP Prefixes
- IP Packets are routed to these prefixes by using their IP destination address
- Peering is exchanging traffic between similar sized networks
 - Peering in the past was mainly done to reduce cost
 - Today peering decreases latency and enhances control
 - Internet Exchanges are places where many networks meet and can peer with each other



Networking Basics 04a - User Datagram Protocol (UDP)

Wolfgang Tremmel academy@de-cix.net

目前) 医肌脂肪的药

Where networks meet

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01 - Networks, Packets, and Protocols

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Internet Model IP / Internet Layer

- Data units are called "Packets"
- Provides source to destination transport
 - For this we need addresses
- Examples:
 - IPv4
 - IPv6



Layer	Nam
5	Applica
4	Transp
3	Interr
2	Linl
1	Physi



Internet Model **Transport Layer**

- May provide flow control, reliability, congestion avoidance
- Examples:
 - TCP (flow control, reliability, congestion avoidance)
 - UDP (none of the above)
- Also may contain information about the next layer up



Layer	Nam
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Encapsulation Packets inside packets

- Encapsulation is like Russian dolls
- IP Packets have a payload
- This payload is usually UDP or TCP (there are others as well)
- So we have an UDP packet inside an IP packet





Attribution: Fanghong. derivative work: Greyhood https://commons.wikimedia.org/wiki/File:Matryoshka_transparent.png



IPv4 Header "Legacy" IP

- Starts with version and length
- Total length of packet
- Important: Time to live (TTL)
- Protocol: Type of payload
 - TCP = 6, UDP = 17
- Source / Destination address 32 bits

Options (optional) **DE CIX**



IPv6 Header Looks simpler, yes?

- Starts with version and some labels
- Payload length in bytes (0-65535
- Next Header you can chain more headers
 - replaces protocol field, same values
- Hop Limit replaces TTL





be	ls
5)	
or	Α

Byte	0	1	2	3		
0	Versio	Version = 6 / Traffic Class / Flow Label				
4	Payloac in b	I Length ytes	Next Header	Hop Lin		
8		Source IPv6 Address				
12						
16						
20						
24						
28		Destination IDv6 Address				
32		Destination IPvo Address				
36						



Next header: Transport layer header TCP, UDP, and more

- We start with the "easiest" protocol
- UDP
 - User Datagram Protocol
 - Protocol ID is 17
 - Introduced in 1980
- Lets have a look at the header





UDP Header

- 4 fields, each of them 16 bits
- Length: UDP header + UDP payload
- Checksum
 - Optional for IPv4, required for IPv6
 - IP header + UDP header are covered
- Source Port
 - Optional, zero if not used
- Destination Port number



required



Port number



Port number









Port numbers In reality...

- Of course we have not a building
- We have a computer system
- But we have port numbers
- Behind each port sits a piece of software
 - On some systems this software is called a "daemon"





UDP - Connectionless communication Why is it called connectionless?

- The sender does not know if and when the packet has been received
- There may be an answer, but there also may be not
- If there is an answer, the sender knows the packet got through
- If there is no answer
 - Either the packet did not get through
 - Or the answer did not get through



UDP packet processing Security issues ahead!

- A UDP packet is received by a system
- It is delivered to the software matching the destination port number
- If a response has to be sent, it is sent to back to sender
 - Using the source IP as destination of the response
 - The source-port becomes the destination port of the response
- Can you see a security problem in that?






UDP - what it is used for



NTP - Network Time Protocol Synchronizing clocks over the Internet

- <u>NTP</u> is a protocol to synchronize computer clocks using the Internet
- Systems send and receive UDP packets on port 123
 - Packets contain a 32-Bit number for seconds and a 32-bit number for fractional seconds
 - Epoch (start) is 1st of January, 1900
 - Rollover will be on 7th of February, 2036





DNS - Domain Name Service The phonebook of the Internet

- DNS translates names (like "<u>www.de-cix.net</u>") to IP addresses
- DNS is so complex and widely used, it deserves a webinar on its own
- Roughly explained
 - A system sends a name to a name server via UDP
 - The name server sends an UDP packet back containing the IP address where the name is hosted







DHCP - Dynamic Host Configuration Protocol This is how your PC gets an IP address at home

- If you connect a computer to a network it needs an IP address
- DHCP takes care it gets one, and more

 - Your computer sends out a DHCP request via UDP broadcast to port 67 A DHCP server replies via UDP and assigns
 - an IP address
 - the default gateway



a nameserver (where to send DNS requests to)



UDP and network security



UDP normal communication Request and answer







UDP as attack tool Faked request and misdirected answer













A real world example Memcached

- <u>memcached</u> is a software to cache objects in RAM for fast retrieval
- Attack method:
 - tell an unsecured installation of *memcached* to store an object
 - that object
 - this gives you an <u>amplification factor</u> of up to 51000





send an UDP packet to that installation with a **faked source IP** to retrieve



A real world example NTP - Network Time Protocol (2010)

- NTP is a protocol to synchronize computer clocks using the Internet
- The "monlist" command, sent via UDP to an NTP server returns the list of the last 600 hosts who have connected to that server
 - If sent from a faked IP source address, this list is sent via UDP to the faked source
- Solution: "monlist" command was removed from the software





Conclusion



Conclusion UDP - User Datagram Protocol

- UDP is a connectionless protocol on the transport layer
- UDP packets are also called "datagrams"
 - the UDP header contains a source and a destination port number
- If misconfigured, UDP services can be used for network attacks
- More and more services which relied on UDP are moved to TCP
 - But TCP is the topic of the next episode



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